

A game between bounded rationals using Genetic Algorithm

Does it pay to be “more rational”?

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Bounded rationality

Why economists have long expressed dissatisfaction with the full rationality assumption that is so pervasive in economic theory?

- Most economic agents are not in fact maximizers
- Many maximization procedures are very difficult to carry out in practice
- Experiments show that people often fail to conform to basic conclusions of rational decision theory
- Results of rational analysis sometimes seem unreasonable even on the basis of simple introspection.

Does it pay to be “more rational”?

Not all the time!

Under full-rationality: There is no issue as to who is more rational

Under bounded rationality: It is an unavoidable question

Who is more rational?



Abe

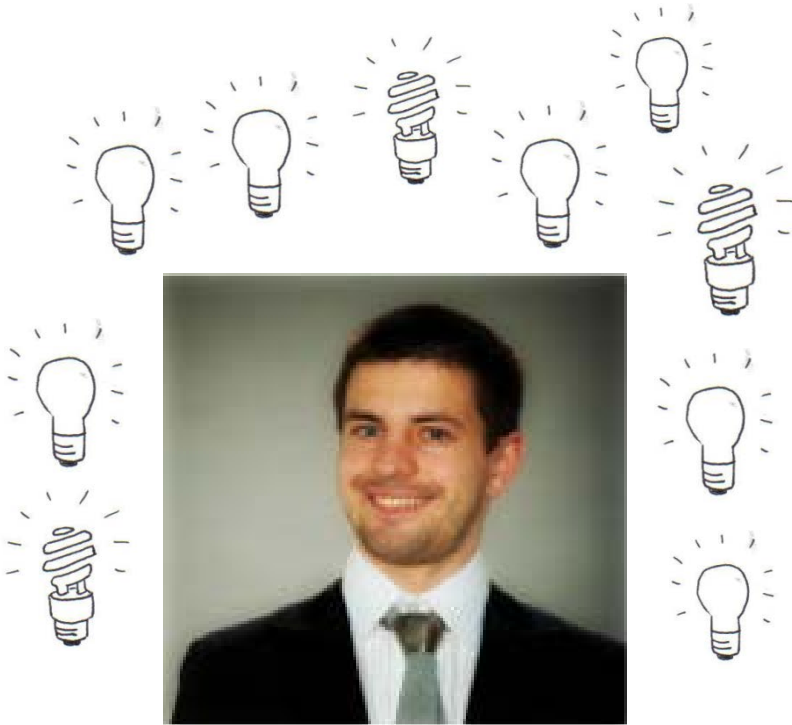
93% rational



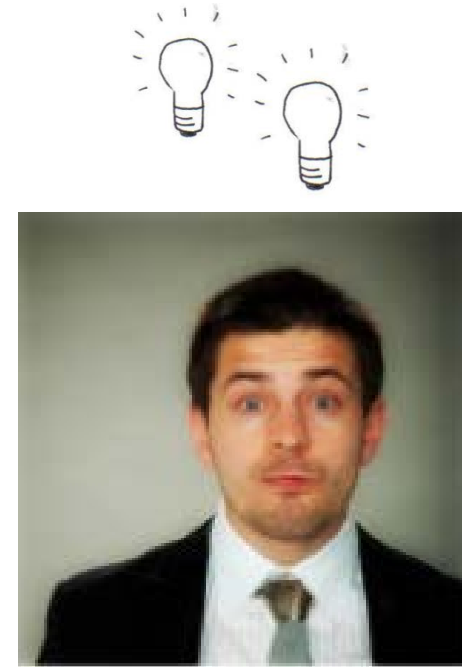
Bob

65% rational

Who is more rational?



Abe



Bob

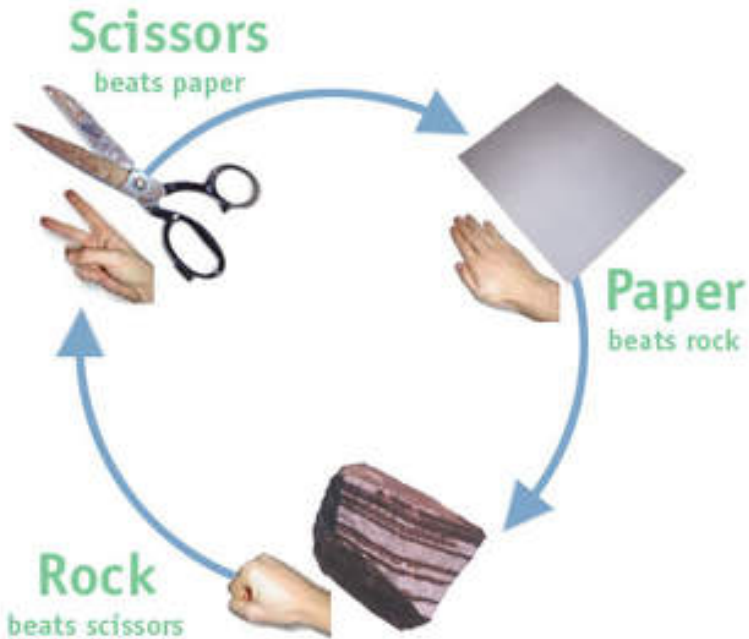
- Can more information really lead to higher payoff?

Assumptions:

- No constraint on computability capacity of individuals (i.e. no information overload or infoxication)
- Marginal utility from additional info is positive and constant
- Each information is distinct and has the same quality (i.e. no info is superior than the other)
- Information sets are tight (i.e. there is no superfluous info in the information set of agents)

All other things constant, we want to see only the effect of the size of the information endowment of an individual on her payoff.

Rock-Paper-Scissors Game



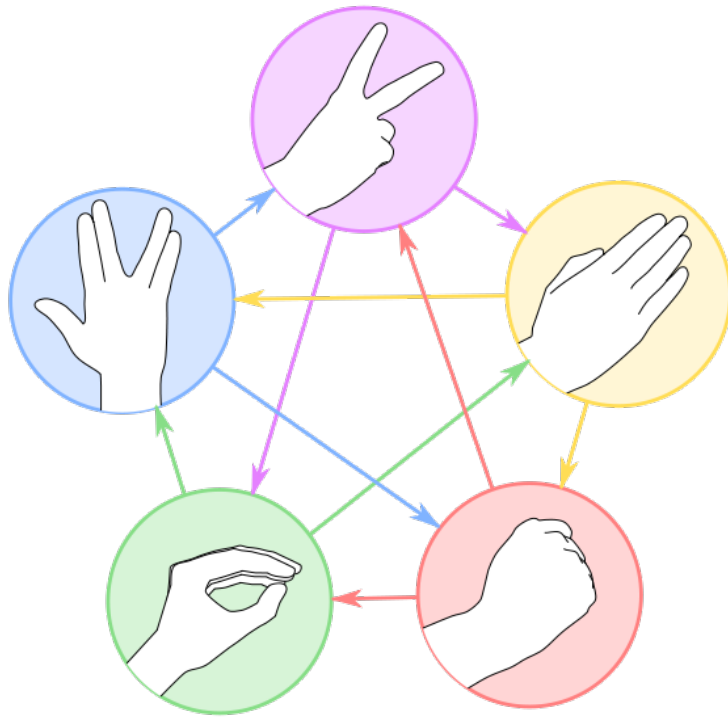
Player 2

R = 1/3 P = 1/3 S = 1/3

Player 1 R = 1/3	0, 0	-1, 1	1, -1
P = 1/3	1, -1	0, 0	-1, 1
S = 1/3	-1, 1	1, -1	0, 0

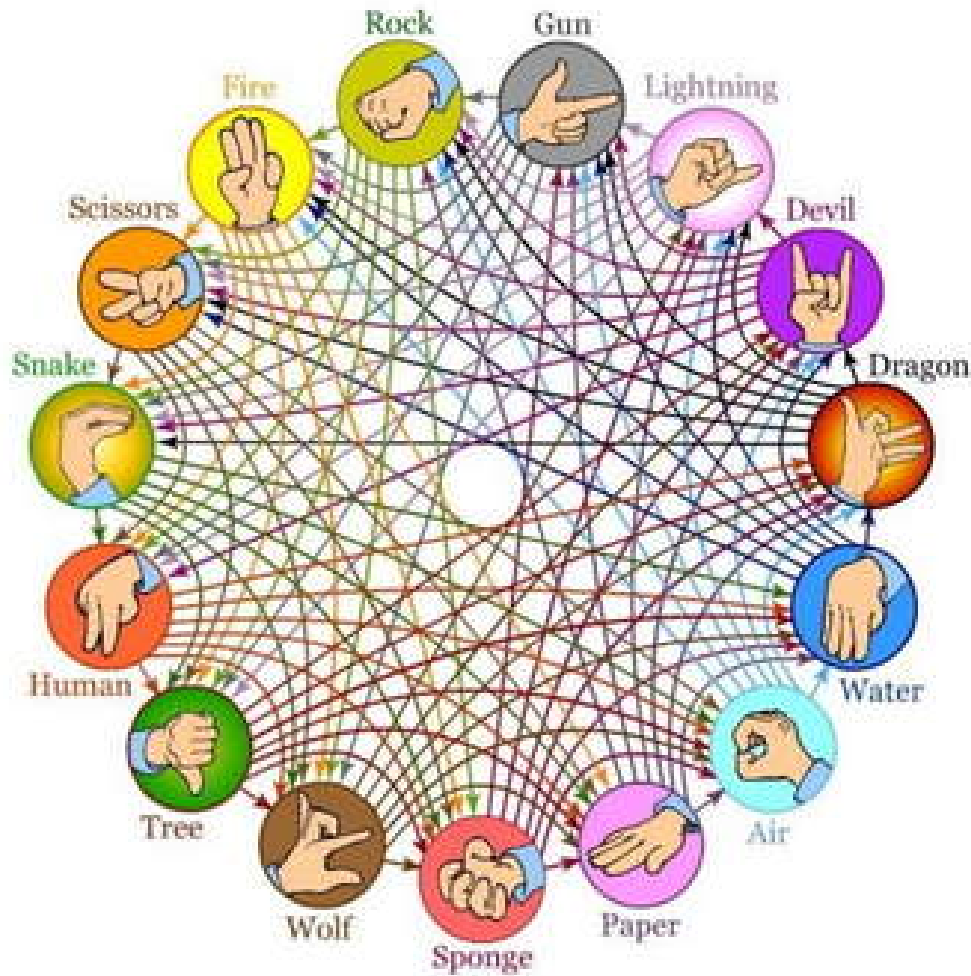
Nash Equilibrium: $[(1/3, 1/3, 1/3), (1/3, 1/3, 1/3)]$

Rock-paper-scissors-lizard-spock



Each option:

- (i) can defeat two other options and
- (ii) can be defeated by another two



For $n=15$, each option can defeat 7 and can be defeated by 7 other

Literature:

- Mobilia, M (2010), “Oscillatory dynamics in rock–paper–scissors games with mutations”, *Journal of Theoretical Biology*.
- Reichenbach, T. et al (2007), “Mobility promotes and jeopardizes biodiversity in rock–paper–scissors games”, *Nature*.
- Szolnoki, A. and G. Szabó (2004), “Phase transitions for rock-scissors-paper game on different networks”, *Physical Review E*.
- Jiang , L et al. (2011), “Effects of competition on pattern formation in the rock-paper-scissors game”. *Physical Review E*.
- Miller, John (1996), “The coevolution of automata in the repeated prisoner’s dilemma”, *Journal of Economic Behavior & Organization*.
- Ali, F and Nakao Z. (2000), “Playing the Rock-Paper-Scissors game using Genetic Algorithm”, *IEEE*.
- Gilboa, I and D. Samet (1989), “Bounded Rationality vs. Unbounded Rationality: The Tyranny of the Weak” , *Games and Economic Behavior*.

A. Both are unbounded:

Player 1

Rock	Lizard
Paper	Spock
Scissors	

VS

Player 2

Rock	Lizard
Paper	Spock
Scissors	

Five-information
universe

B. One is unbounded, the other bounded:

Player 1

Rock	Lizard
Paper	Spock
Scissors	

VS

Player 2

Rock	Lizard
Paper	

C. Both are bounded:

Player 1

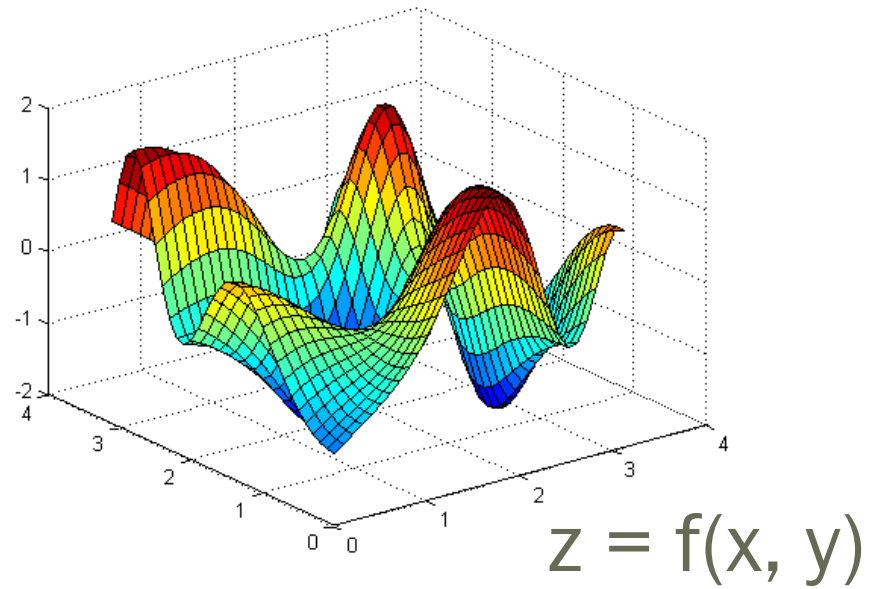
Rock
Spock

VS

Player 2

Rock	Lizard
Paper	

Optimization:



$$\Pi_x = f(x_1, x_2, x_3; y_1, y_2, y_3)$$

$$\Pi_x = f(R, P, R; S, S, S)$$

$$\Pi_x = 2 \quad \Pi_y = 1$$

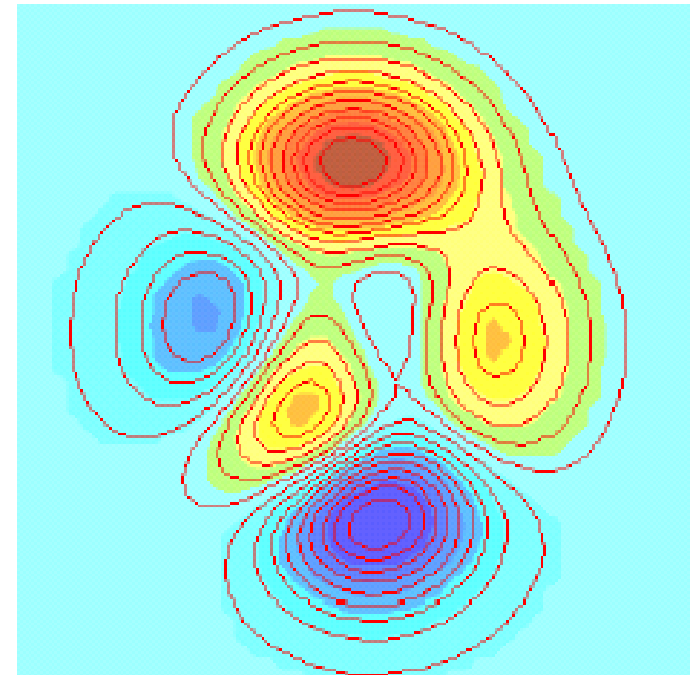
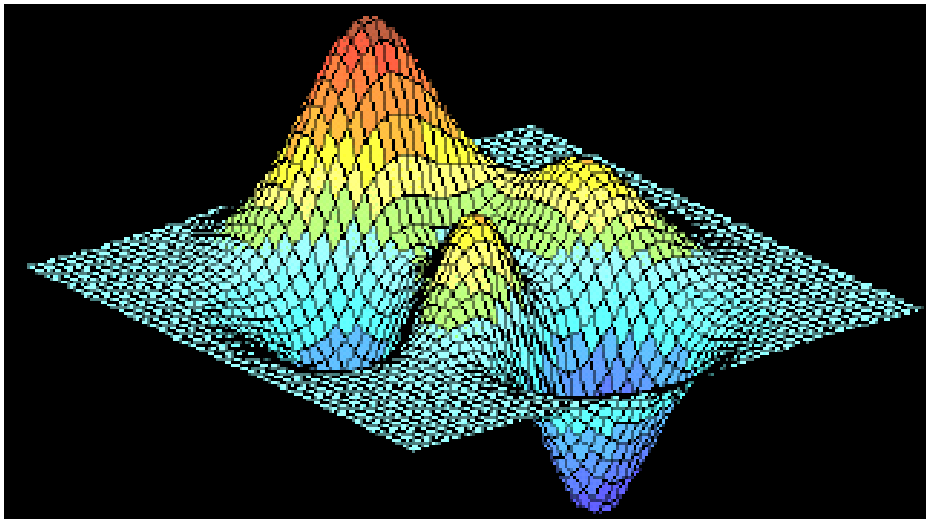
f is qualitative

Genetic Algorithm (GA)

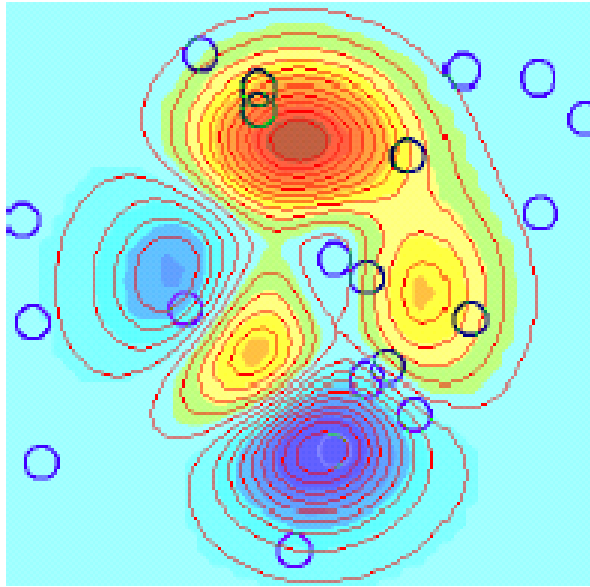
- a method for solving optimization problems based on a natural selection process that mimics biological evolution.
- the algorithm randomly selects individuals from the current population and uses them as “parents” to produce “children” for the next generation.
- over successive generations, the population evolves toward an optimal solution.

Typical optimization problem using GA

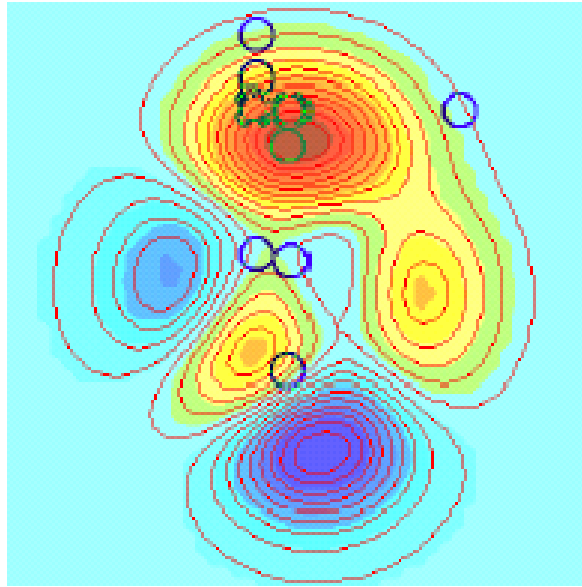
- $z = f(x, y) = 3 \cdot (1-x)^2 \cdot \exp(-(x^2) - (y+1)^2) - 10 \cdot (x/5 - x^3 - y^5) \cdot \exp(-x^2 - y^2) - 1/3 \cdot \exp(-(x+1)^2 - y^2)$.



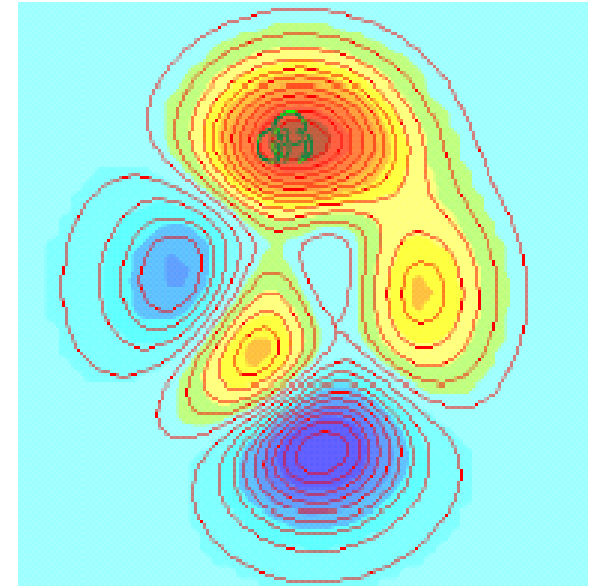
- GA process:



Initial population



5th generation



10th generation

Simplifying analogy

- Chromosome \longleftrightarrow Strategy (7-bit strategy)
- Genes \longleftrightarrow {R, P, S, etc.}
- Genotype \longleftrightarrow (R, P, S, S, P, S, S)
- Phenotype \longleftrightarrow {1,0,-1}

GA Mechanism

1. Initialization: generate an initial population at random (i.e. $P(i)=P(1)$, first generation).

2. Iterate:

- a) evaluate the fitness of individuals in $P(i)$
- b) select parents from $P(i)$ based on their fitness
- c) generate offspring from the parents using crossover & mutation to form $P(i+1)$
- d) go back to (a) with $P(i+1)$

3. Continue until some stopping criteria is satisfied.

Report statistics!

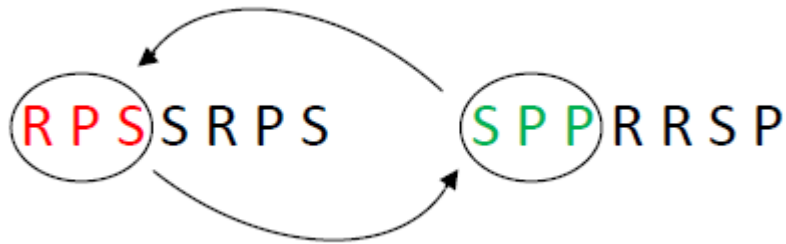
Reproduction Methods

1. Cross-over

1st Generation (parents)

RPSSRPS
Parent 1

SPPRRSP
Parent 2



2nd Generation (offspring)

SPPSRPS

RPSRRSP

2. Mutation

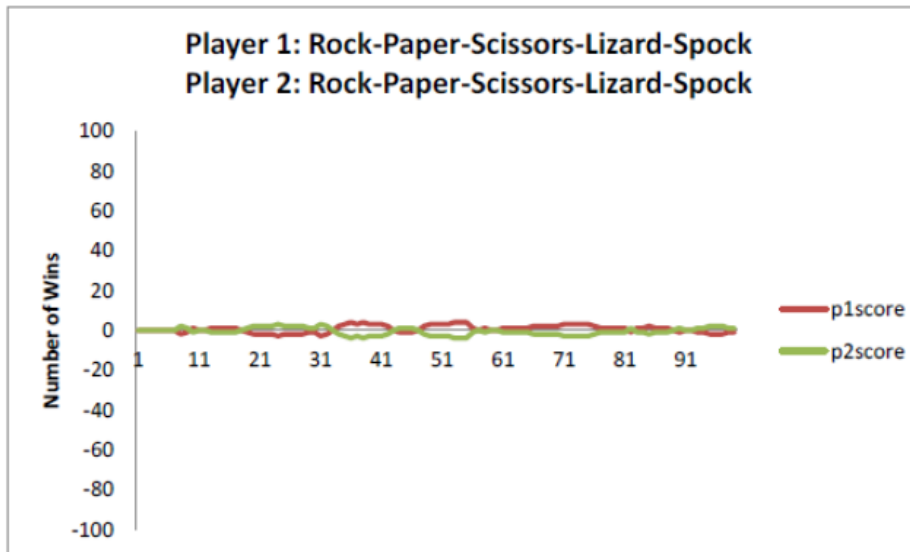
RPSSRPS



RRSSRPP

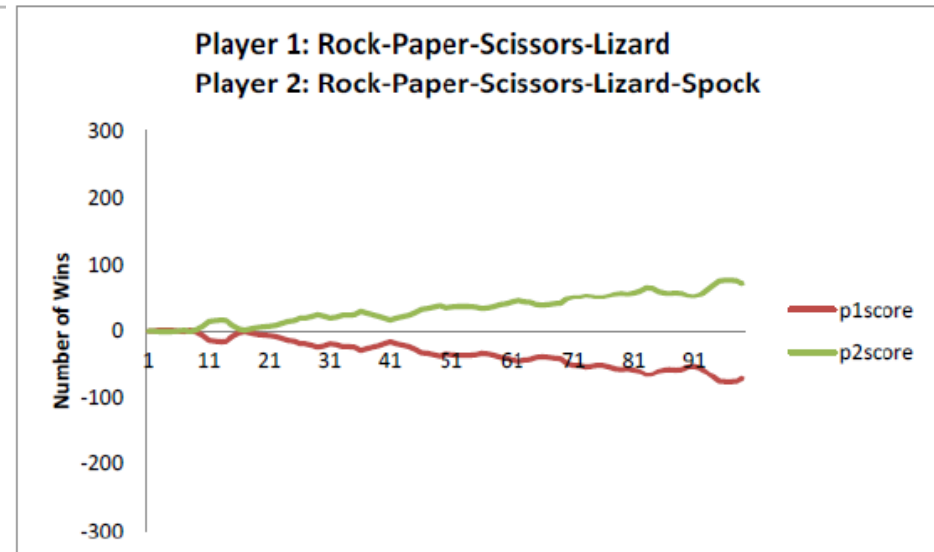
Simulation Results:

A. Both are undounded:



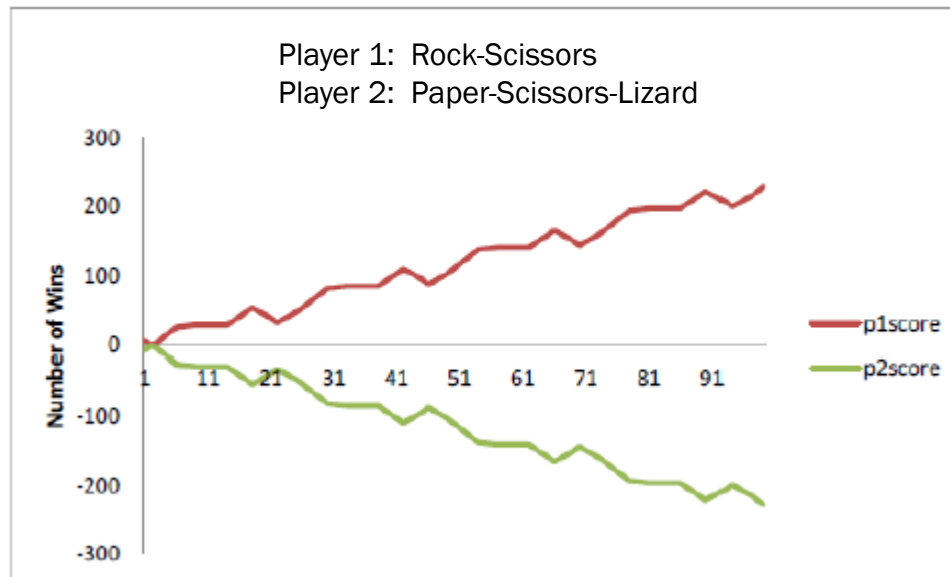
On the average, the result is a draw.

B. One is bounded,
the other is unbounded:

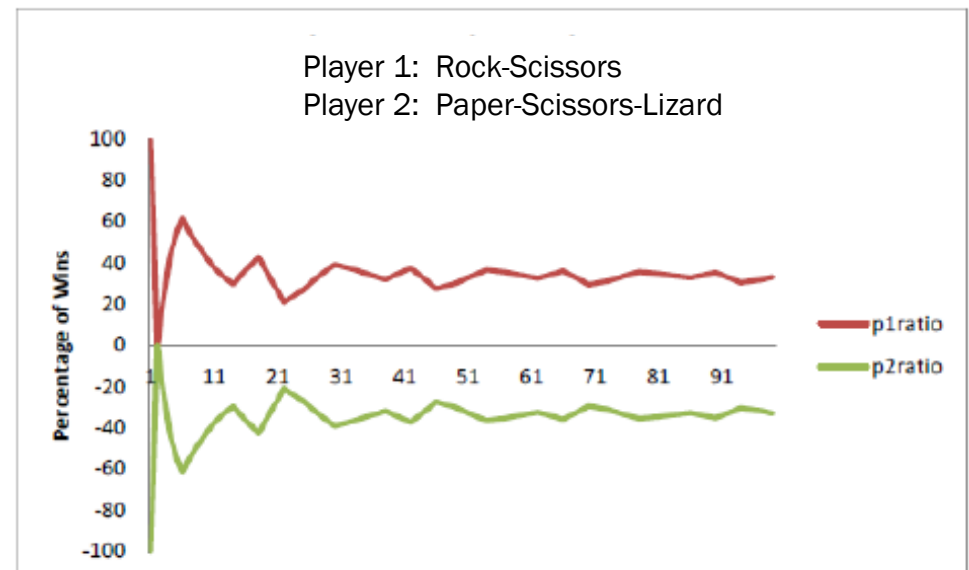


The bounded fellow is almost
always defeated.

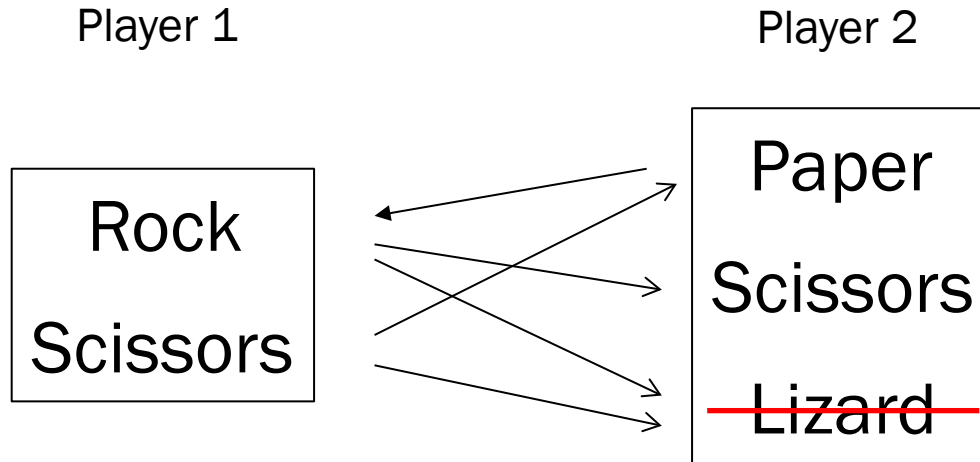
C. Both are bounded:



Player 1 who has less info is getting higher payoff !



Elimination of weakly dominated strategies:



Player 2

	Paper	Scissors	Lizard
Player 1 Rock	-1, 1	1, -1	1, -1
Scissors	1, -1	0, 0	1, -1

Lemma: In a zero-sum game with a full set of actions Ω for both players, any arbitrary deletion of an action $a_y \in \Omega$ by player y makes a certain action $a_x \in \Omega$ by player x weakly dominated.

Proof:

Let a_y be the action deleted from the roster of actions of player y . Since there are $(n - 1)/2$ actions of x that beats a_y , these actions a_x^- will now have negative expected payoffs i.e.

$$a_x^- = \left\{ a_x \in \Omega_x \mid \sum_{y=1}^{n-1} u_x(a_x, a_y) P_y < 0 \right\}.$$

To show that one of these a_x^- is weakly dominated, pick an action a'_x such that it can beat the same set of y actions that a a_x^- can defeat. There are two of these a_x^- that can fulfill this: One that defeats a'_x and another that is defeated by a'_x . By choosing the latter, we now have an $a_x^* \in a_x^-$ whose $u_x(a'_x, \tilde{a}_y) \geq u_x(a_x^*, \tilde{a}_y)$ for all $\tilde{a}_y \in \Omega \setminus a_y$. The action a_x^* is therefore weakly dominated brought about by the removal of action a_y . ■

Proposition 1: A player (fully-rational) who employs all the possible actions in a finite universal set Ω always obtains a higher expected payoff than a player (bounded rational) who only employs a set of actions A , where $A \subset \Omega$.

Proposition 2: For any action set $A \subset \Omega$ that is available to player x , there exists an action set $B \subset \Omega$ available to player y (where the number of elements in B is less than the number of elements in A) such that the expected payoff it gives to y is higher than what x can obtain from employing A .



VS



The fellow with less options wins around 65 percent of the time!



The End

Thank you!